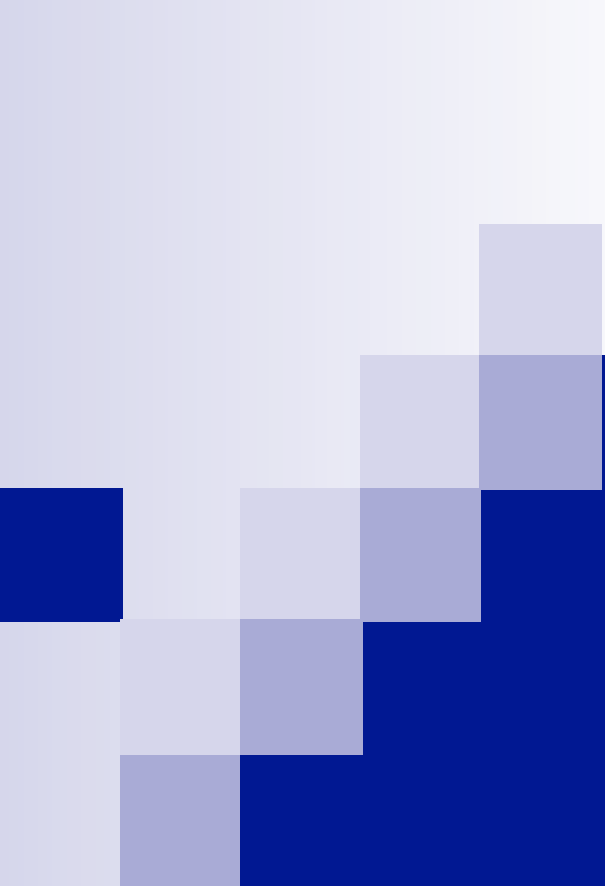




Distributed Systems Lecture 1

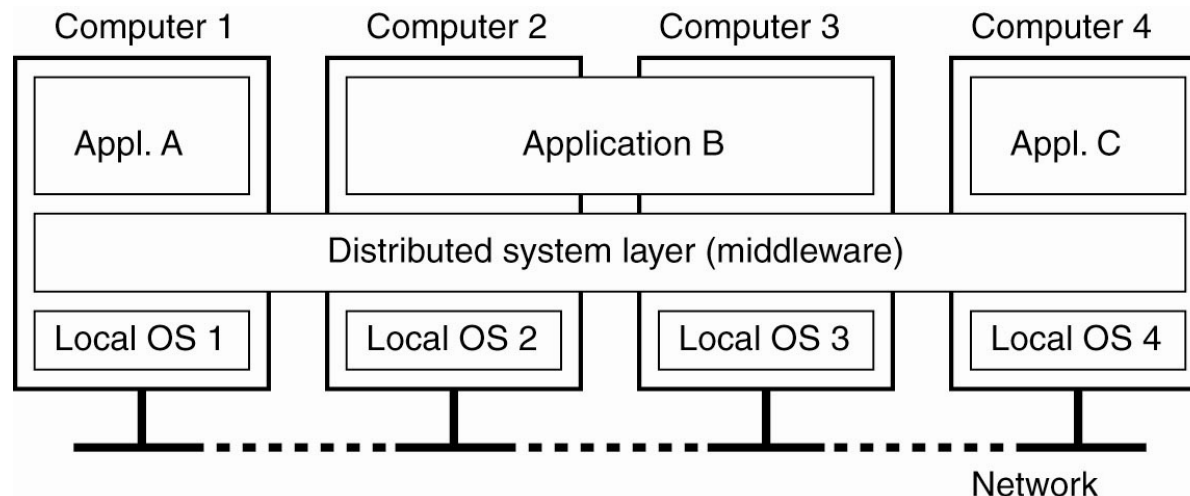
Professor Louise E. Moser
Spring 2014

Course Administrative Details

- **Lectures:** T Th 11am-12:15pm Phelps 2524
- **Discussion Sections:** Th 5-5:50pm Phelps 1444
F 10-10:50am Phelps 1440
- **Office Hours:**
 - Prof Louise Moser: T Th 2-3:30pm HFH 2164
 - TA Miguel Lastras Montano: M W 3-4pm, Th 4-5pm Phelps 1435
 - TA Harichandan Pulagam: M W 1-2pm, F 11-12noon Phelps 1435
- **Text:** Distributed Systems: Principles and Paradigms
A. S. Tanenbaum and M. van Steen, Prentice Hall, Second Edition
- **Midterm:** Tue, May 6, 11-12:15pm, **Final:** Wed, June 11, 12-3pm
- **Grading:** Midterm 15%, Final 35%, Assignments 35%
Class Attendance / Participation 15%
(Lectures and Discussion Sections)

Definition of a Distributed System

- A **Distributed System** is a collection of two or more independent computers that appears to its users as a single coherent system
- A distributed system is typically organized as middleware that extends across multiple computers and that offers each application the same interface





Goals of a Distributed System

- Overall Goal: Connecting resources and users
- Additional Goals:
 - Transparency
 - Openness
 - Scalability

Different Forms of Transparency

Transparency	Description
Access	Hide differences in data representations and how a resource is accessed
Location	Hide where a resource is located
Relocation	Hide that a resource may be moved to another location
Migration	Hide that a resource may be moved to another location while it is being used
Replication	Hide that a resource may be available on several distinct computers
Concurrency	Hide that a resource may be shared by several competing users at the same time
Failure	Hide the failure and recovery of a resource

Degree of Transparency

Complete transparency might not be desirable or possible

- Users might be located in different countries - distribution is apparent and might not be something you want to hide
- Complete transparency can result in extra performance costs
 - **Keeping Web caches *exactly* up-to-date with the master copy**
 - **Immediately copying the results of write operations to disk**
- Hiding faults in hardware and software might not be theoretically or practically possible
 - **You cannot distinguish a slow computer from one that has failed**
 - **You cannot be sure that a computer has actually performed an operation before it crashed**

Openness

- An **open** distributed system is one that interacts with services from other open systems and requires
 - Well-defined interfaces and protocols
 - Portability of applications
 - Interoperability of systems
- An open distributed system is one that is independent of the heterogeneity of the underlying environment
 - Languages
 - Operating systems
 - Hardware

Implementing Openness

- **Policies** specified by the applications and the users
 - What level of consistency do we require for client-cached data?
 - Which operations do we allow downloaded code to perform?
 - Which QoS requirements do we adjust for different applications?
 - What level of security/privacy do we require for communication?

- **Mechanisms** implemented by the middleware and the operating system
 - Allow dynamic setting of caching policies, preferably per cached item
 - Support different levels of trust for downloaded code
 - Provide adjustable QoS parameters per data stream
 - Offer different encryption algorithms

Scalability

■ Limitations of scalability due to centralization

Concept	Example
Centralized services	A single server for all users
Centralized data	A single on-line telephone book
Centralized algorithms	Doing routing based on complete information

Scalability

- Characteristics of **decentralized** algorithms
 - No machine has complete information about the state of the distributed system
 - Machines make decisions based on local information
 - No implicit assumption that a global clock exists
 - Failure of one machine does not affect the operation of the algorithm

Scalability Metrics

- Many developers of distributed systems use the term “scalable” without making clear *how* their systems scale
- At least three scalability metrics exist
 - **Size scalability** – number of users or processes or machines
 - **Geographical scalability** – maximum distance between nodes
 - **Administrative scalability** – number of administrative domains

Scaling Techniques

- **Distribution:** Partition the data and computations across multiple machines
 - Move computations to clients (Java applets)
 - Decentralized naming services (DNS)
 - Decentralized information systems (WWW)
- **Caching:** Allow client processes to access local copies
 - Web caches (browser or proxy)
 - File caching (client or server)
- **Replication:** Make copies of the data available on different machines (mainly for fault tolerance)
 - Replicated file servers
 - Replicated databases
 - Mirrored Web sites

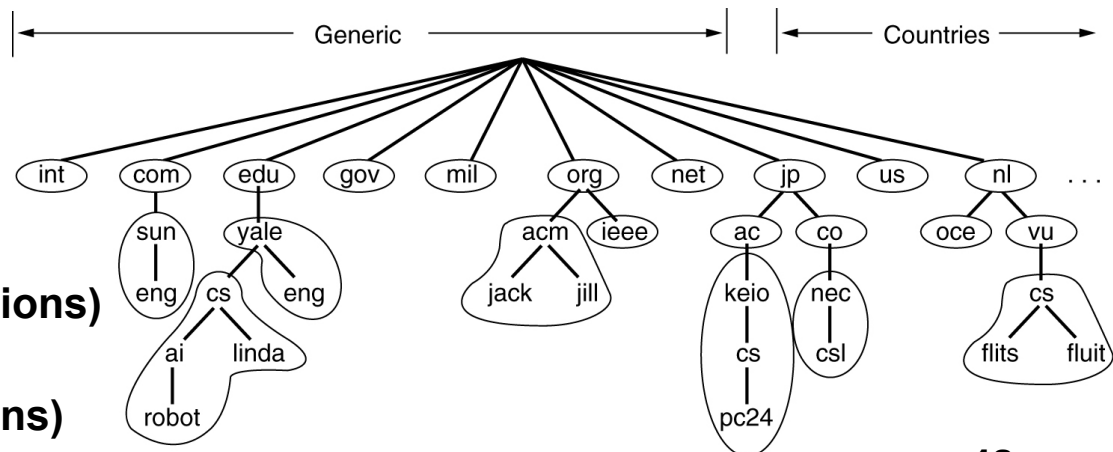
Example Scaling Technique

- **Domain Name System (DNS)** in the Internet is hierarchically structured into several hundred top-level domains

- Domains correspond to organizational boundaries, *not* physical boundaries
- Domains are partitioned hierarchically in a tree structure into subdomains, etc.
- To create a new domain, need permission from the manager above in the tree
- Domain is named by a path, e.g., ece.ucsb.edu

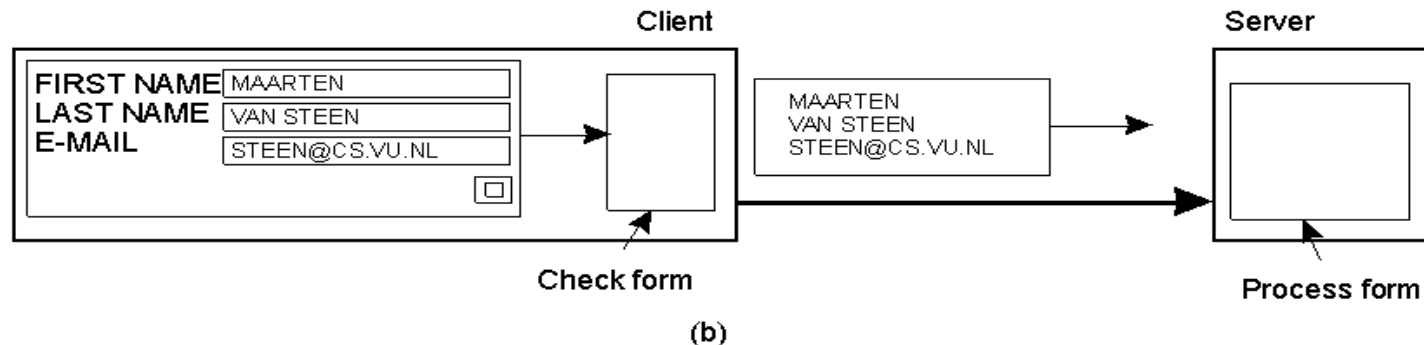
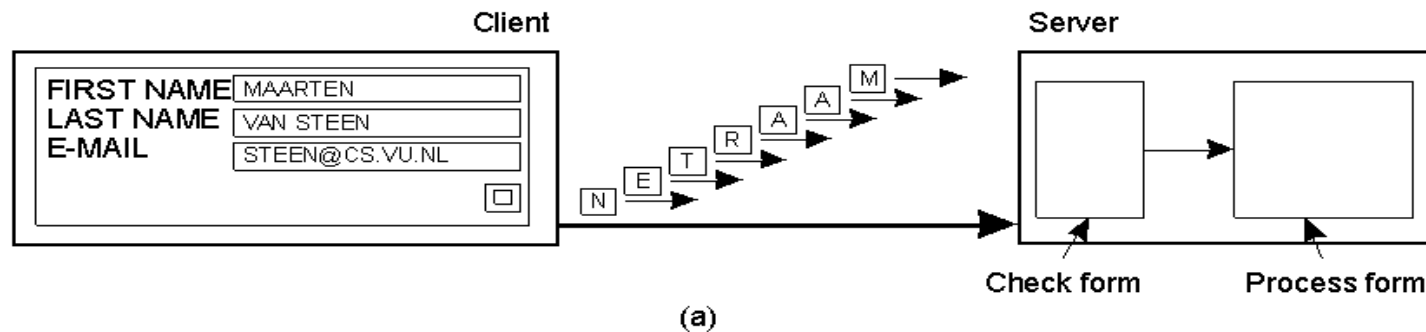
- **Generic DNS names**

- **com** (commercial)
- **edu** (educational)
- **gov** (US government)
- **mil** (US military)
- **int** (international organizations)
- **net** (network providers)
- **org** (non-profit organizations)



Example Scaling Technique

- (a) Server checks the form for each client, or
- (b) Client checks the form, which is more scalable because server can then serve more clients



Pitfalls in Developing Distributed Systems

- **False assumptions** made by a first-time developer of a distributed system
 - Network is homogeneous
 - Network is secure
 - Network is reliable
 - Topology does not change
 - Latency is zero
 - Bandwidth is infinite
 - Transport cost is zero
 - There is one administrator

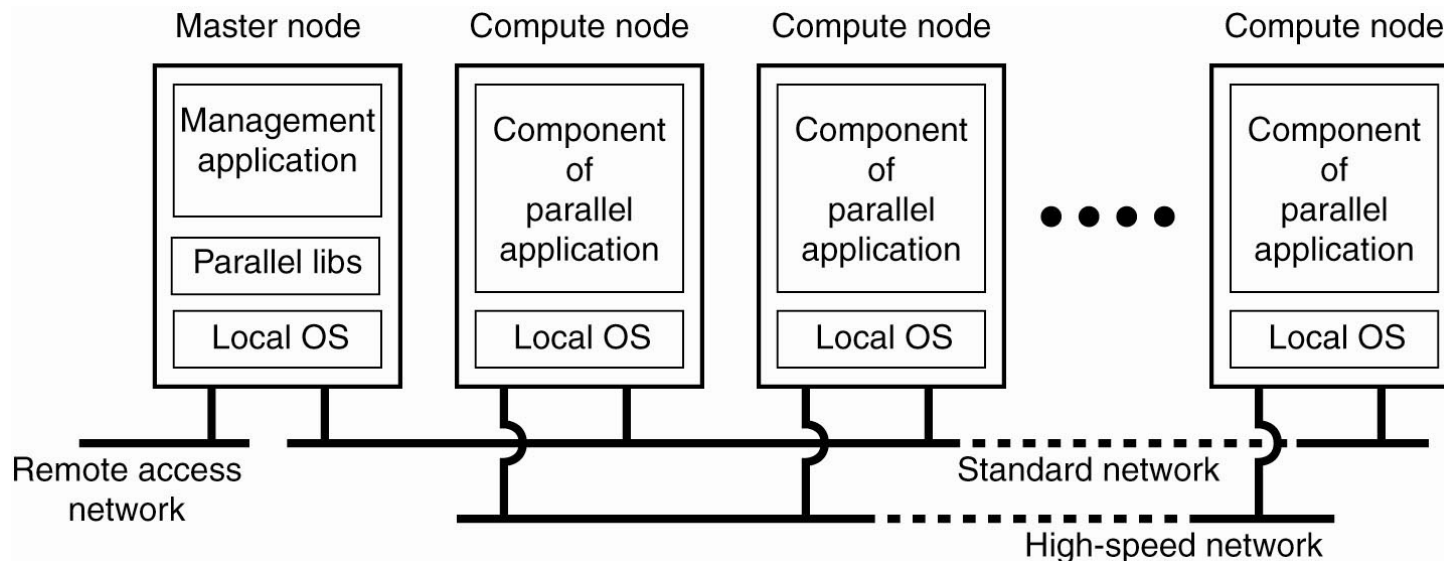


Types of Distributed Systems

- **Distributed Computing Systems**
 - **Cluster Computing Systems**
 - **Grid Computing Systems**
- **Distributed Information Systems**
 - **Transaction Processing Systems**
 - **Enterprise Application Integration (EAI)**
- **Pervasive Distributed Systems**
 - **e-Healthcare Systems**
 - **Sensor Networks**

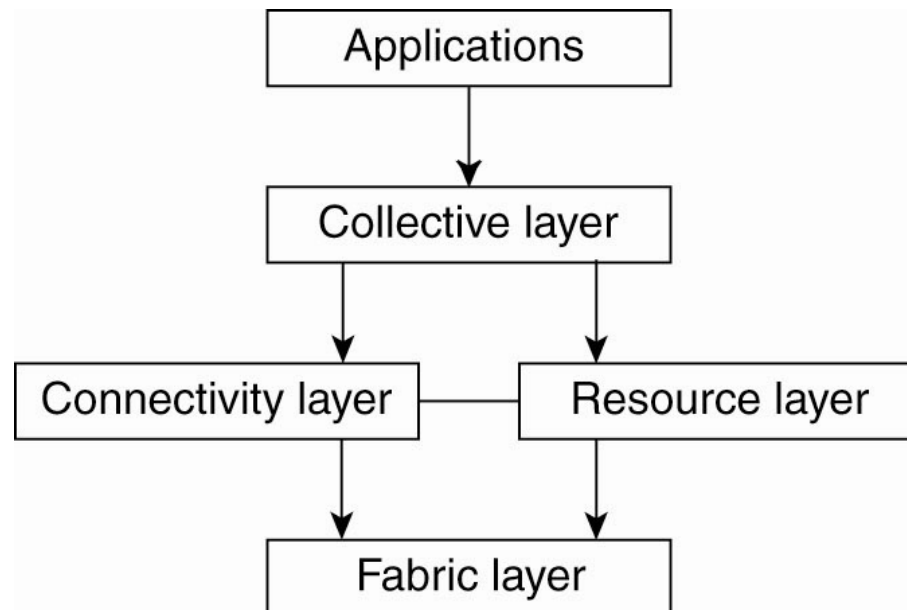
Cluster Computing Systems

- A **cluster** is a collection of workstations or PCs connected by a local-area network (LAN)
- Each node runs the same operating system (OS)
- Example: Master/slave system where the slaves (compute nodes) operate in parallel



Grid Computing Systems

- A **grid** is a federation of computing systems
 - Each system might be in a different administrative domain
 - Each system might use different software, hardware, network technology
- A layered architecture for a grid computing system



Transaction Processing Systems

- A **transaction** is a group of program statements that are intended to be executed together (i.e., atomically)
- Standard primitives for transactions

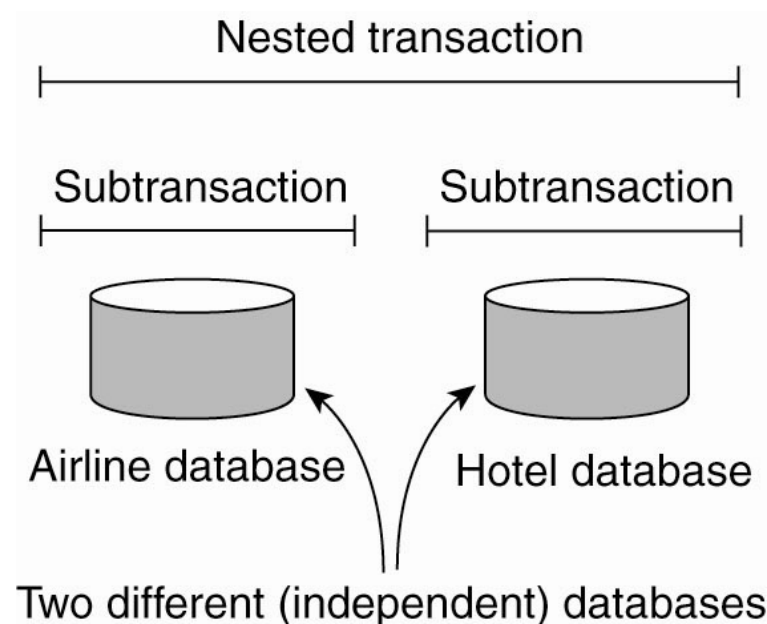
Primitive	Description
BEGIN_TRANSACTION	Mark the start of a transaction
END_TRANSACTION	Terminate the transaction and try to commit
ABORT_TRANSACTION	Kill the transaction and restore the old values
READ	Read data from a file, a table, or otherwise
WRITE	Write data to a file, a table, or otherwise

Transaction Processing Systems

- Traditionally, transactions operate on a single local database
- Such transactions are expected to satisfy the following **ACID** properties:
 - **Atomic:** To the outside world, the transaction happens indivisibly, i.e., all of the operations of the transaction are performed, or none of them is performed
 - **Consistent:** The transaction does not violate system invariants
 - **Isolated:** Concurrent transactions do not interfere with each other; they can be serialized
 - **Durable:** Once a transaction commits, the changes are made permanent, by writing them to disk
- Note: It is very difficult to guarantee the ACID properties for distributed transactions that operate on multiple databases and that span multiple enterprises

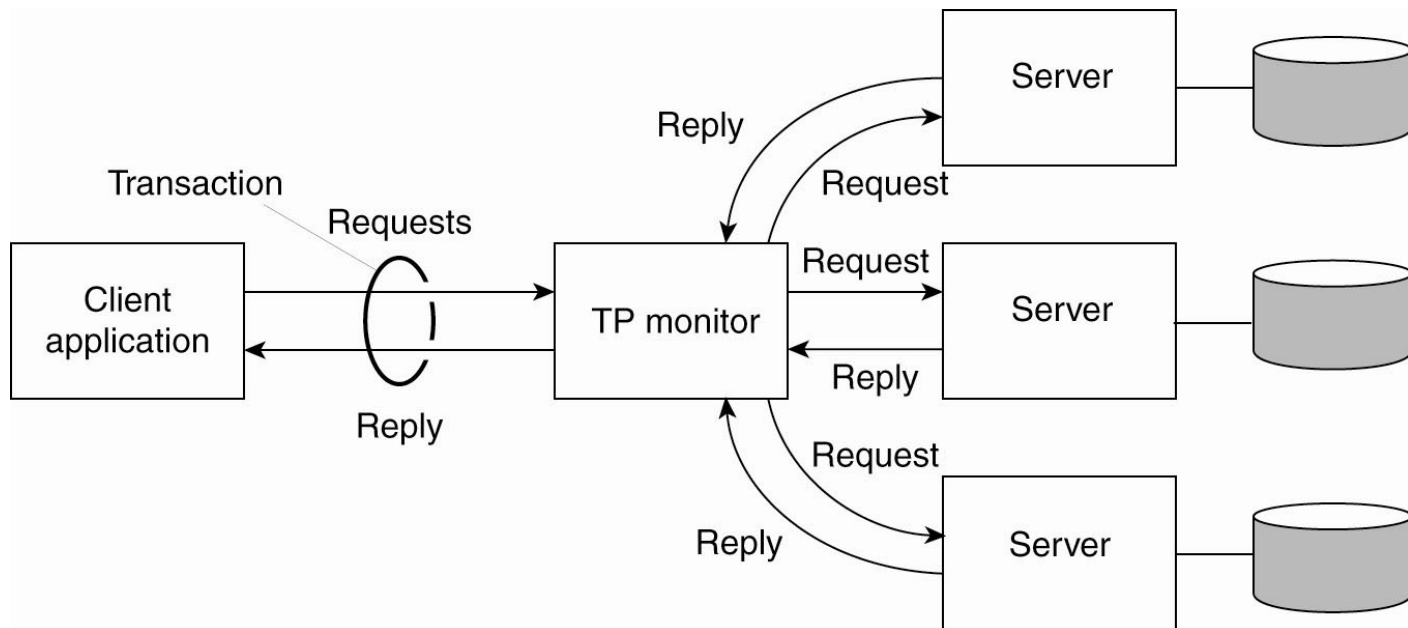
Transaction Processing Systems

- **Nested transactions** with **subtransactions**
- Subtransactions typically operate in parallel on different databases
- Durability applies only to the top-level transactions



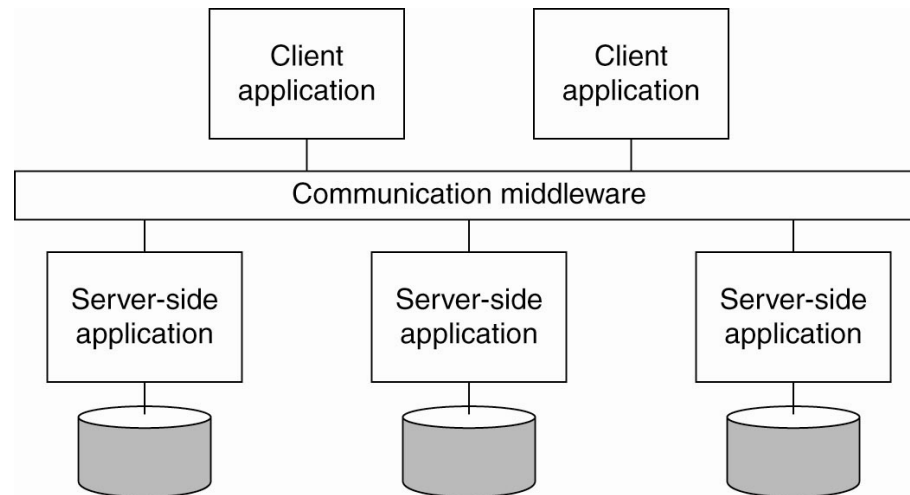
Transaction Processing Systems

- **Transaction Processing (TP) monitor** allows an application to access multiple servers/databases, using the transaction model



Enterprise Application Integration

- In **Enterprise Application Integration (EAI)**, a distributed application is separated into independent components
 - In particular, processing components are separated from database components
- These components need to be integrated and to communicate with one another
 - **Middleware facilitates this integration and communication in EAI**



Pervasive Distributed Systems

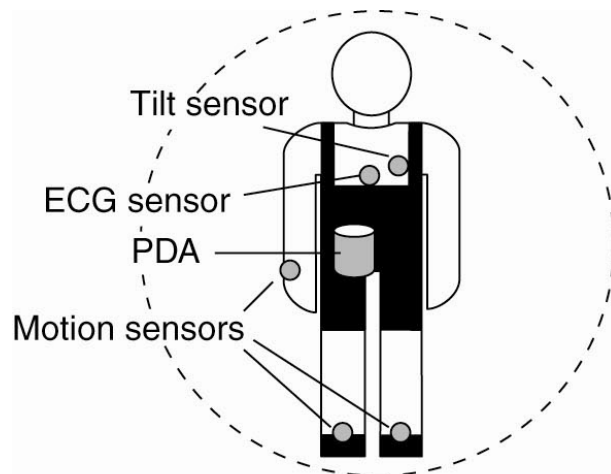
- A **pervasive distributed system** typically comprises small, battery-powered, mobile computing devices with wireless connections
 - Part of our surroundings, lacks human administrative control
 - Devices must automatically discover their environment and other devices within it
- Pervasive distributed systems requirements:
 - Adapt to contextual changes
 - Enable ad hoc composition
 - Recognize sharing as the default

Example: e-Healthcare Systems

- Questions to be addressed in an e-healthcare system
 - Where and how should patient data from monitoring devices be stored?
 - How can we prevent loss of crucial patient data?
 - What infrastructure is needed to generate and propagate alerts to the patient, caregivers, healthcare professionals?
 - How can physicians provide online feedback?
 - What are the security and privacy issues and how can the proper policies be enforced?
 - How can robustness of the e-healthcare system be realized?

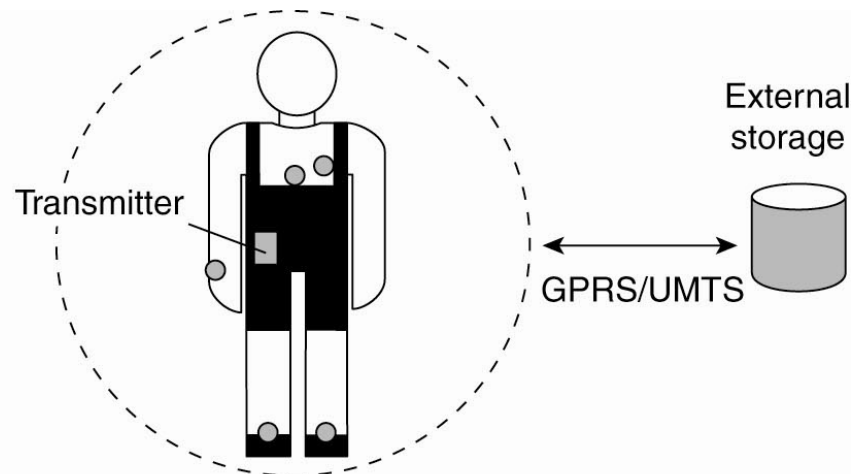
e-Healthcare Systems

- Health monitoring of a patient using
 - (a) A local hub
 - (b) A continuous wireless connection



body-area network

(a)



body-area network

(b)

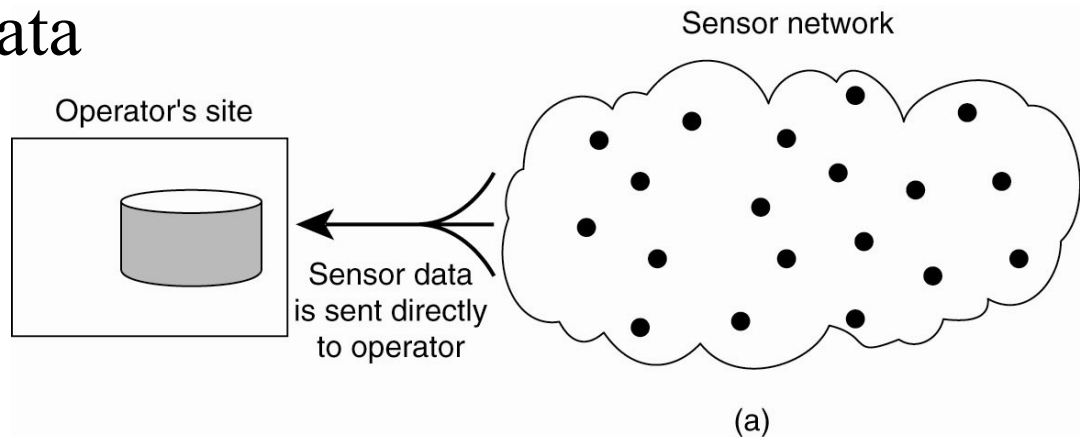
Example: Sensor Networks

- A **sensor network** typically comprises hundreds to thousands of small battery-powered nodes (each with a sensor) that use wireless communication
 - Typically, used for measurement and surveillance
 - Efficiency is a prime design criterion due to limited processing and memory resources, communication capabilities, and power supplies
- Questions concerning data processing within a sensor network
 - How do we set up an efficient means of communication?
 - How do we aggregate and consolidate results?
 - What happens when a network link fails?

Sensor Networks

- A sensor network database with storage and processing of data

(a) Only at the operator's site



(b) Only at the sensors

