

LECTURE #3

- Path loss model;
 - holds for "far-field" of transmit antenna
'Fraunhofer region'

d_f : far-field distance.

$$d_f = \frac{2D^2}{\lambda}$$

D : largest linear dimension of the transmitter antenna aperture.

λ : carrier wavelength.

- The reference distance d_0 must be chosen to be in the far field:

$$\left\{ \begin{array}{l} d_0 \geq d_f = \frac{2D^2}{\lambda} , \text{ and} \\ d_0 \gg D , \text{ and} \\ d_0 \gg \lambda \end{array} \right\}$$

Example 1:

- Find the far-field distance of an antenna with maximum dimension 1 m, and center frequency 900 MHz

Solution: $D \approx 1 \text{ m}$,

$$f_c = 900 \text{ MHz}$$

$$d_f = \frac{2 D^2}{\lambda} = \frac{2 (1)^2}{\left[\frac{3 \times 10^8 \text{ m/s}}{900 \times 10^6} \right]} = 2 \cdot 3 = \underline{6 \text{ meters}}$$

Far field: $\geq 6 \text{ meters}$, and

$\gg 1 \text{ meter}$, and

$\gg \frac{1}{3} \text{ meter}$,

$\therefore \underline{\approx 10 \text{ meters}}$ will fall in the far-field.

Example 2: Assume $P_t = 50 \text{ watts}$

omnidirectional, ^{tx} antenna, $f_c = 900 \text{ MHz}$; omnidirectional receive antenna

- (a) Find received power in dBm at a distance of 100 m from the antenna in free space,

- (b) $P_r (10 \text{ km}) = ?$
in free space

Solution:

$$P_r = \frac{P_t \cdot G_t \cdot G_r \cdot \lambda^2}{(4\pi)^2 \cdot d^2 L}$$

↑ assume 1.

$$= \frac{50 \cdot 1 \cdot 1 \cdot \left[\frac{1}{3} \text{ metres}\right]^2}{(4\pi)^2 \cdot (100)^2}$$

$$= 3.5 \mu W$$

$$P_r \text{ (dBm)} = 10 \log_{10} (3.5 \times 10^{-3} \text{ mW}) = \underline{\underline{-24.5 \text{ dBm}}}$$

(b) Using the path loss model: (free space $n=2$)

$$P_r(10 \text{ km}) = P_r(100 \text{ metres}) + 20 \log_{10} \left[\frac{\overset{100 \text{ meter}}{\underset{100}{(d_0)}}}{\underset{10 \text{ km}}{10000 \text{ metres}}} \right]$$

$$\begin{aligned} \left[P_r(10 \text{ km}) \right. &= P_r(100 \text{ metres}) - 10 \log_{10} \left(\frac{d}{d_0} \right)^2 \left. \right] \\ \text{(dB)} &\quad \text{(dB)} \quad \text{(dB)} \\ &= -24.5 \text{ dBm} - 40 \text{ dB} \\ &= \underline{\underline{-64.5 \text{ dBm}}} \end{aligned}$$

Lognormal shadowing.
 ("large-scale variations")

$$PL(d) = \overline{PL}(d) + X_\sigma$$

$$X_\sigma \sim \text{Gsn}(0, \sigma^2)$$

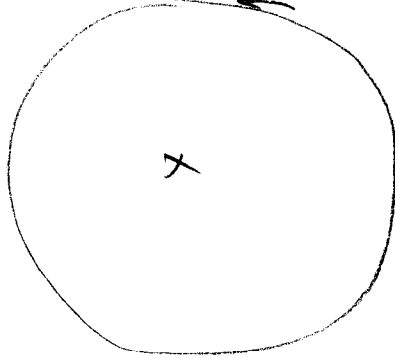


- you'll find that values ~~are~~ (in dB) are normally distributed.

$$P_r = \underbrace{\alpha \cdot P_t \cdot \left(\frac{d}{d_0}\right)^{-n}}_{\substack{\text{path loss} \\ \text{return value} \\ \propto 1/d^\alpha}} \cdot \underbrace{10^{X_\sigma/10}}_{\substack{\text{random var} \\ \text{— lognormally} \\ \text{distributed}}}$$

- Arises from obstacles (trees, buildings) between Tx and Rx

Probability of outage: = P_r (the received signal power falls below a certain threshold),

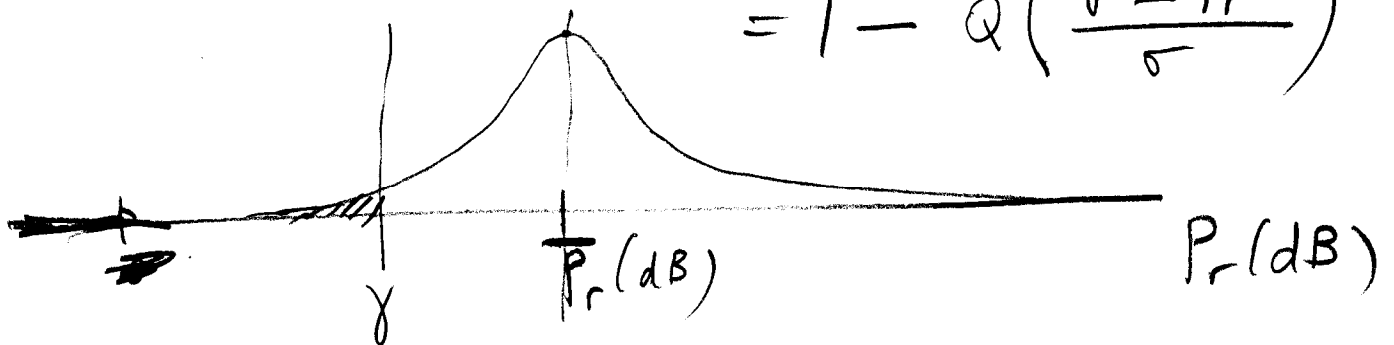


e.g. 5%, 1%.

($\Rightarrow$ call will be dropped)

~~$P_r(P_r)$~~

$$P\{\text{~~signal~~ } P_r < \gamma\} = 1 - P_r(P_r > \gamma) \\ = 1 - Q\left(\frac{\gamma - \bar{P}_r}{\sigma}\right)$$



$\sigma = 8-12$ dB in macrocells.

[Show slides: Measurements from urban and suburban environments.]

Example: Consider a single cell in a cellular system. Design the system such that $P_{out} = 0.01$ (1%) for the worst-case user (on cell boundary).

SYSTEM
PARAMS: $f_c = 900 \text{ MHz}$
Antenna dimensions ≈ 1 meter.

Environment
PARAMS: $\sigma = 8 \text{ dB}$
(of shadowing)
 $n = 4$.

$$d_{max} = 10 \text{ km.}$$

(maximum distance from the B.S.)

$$P_t = 50 \text{ Watts}$$

- Assume omnidirectional antennas.

Solution: First, let's determine the reference distance d_0 .

$$d_f = \frac{2D^2}{\lambda} = \frac{2(1)^2}{\left[\frac{3 \times 10^8 \text{ m/s}}{900 \times 10^6 \text{ /s}} \right]} = 6 \text{ meters}$$

To be in the far field, $\gg 6$ meters,

$$\gg D = 1 \text{ meter}$$

(LOS)

$$\gg \frac{\lambda}{3} \text{ meter}$$

using free space model.

take $d_0 = 10 \text{ meters}$

$$P_r(10 \text{ meters}) = \frac{P_t \cdot G_t \cdot G_r \cdot \lambda^2}{(4\pi)^2 \cdot d^2} = \frac{50 \cdot 1 \cdot 1 \cdot \left[\frac{1}{3}\right]^2}{(4\pi)^2 \cdot 10^2}$$

$$\left(\frac{100^2}{12^2} = 10000 \right)$$

$$= 0.35 \text{ mW}$$

$$[\bar{P}_r(d) = \bar{P}_r(d_0) \cdot \left(\frac{d}{d_0}\right)^{-n}]$$

Then:

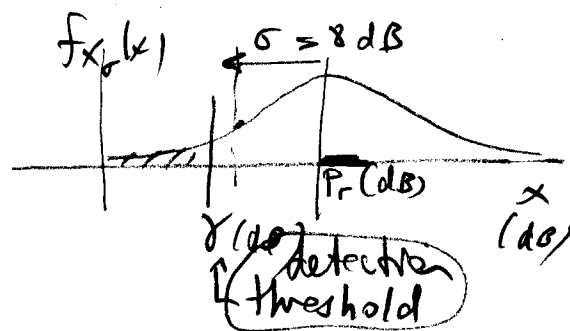
$$\bar{P}_r(10 \text{ km}) = (0.35 \text{ mW}) \cdot \left(\frac{10 \times 10^3}{10}\right)^{-4}$$

$$= 0.35 \times 10^{-12} \text{ mW}$$

↑
average power level at 10 km.

Now, $P_{\text{out}} = 0.01$, which means

we want:



$$P\{\bar{P}_r(d_{\text{max}}) > \gamma\} = 0.99$$

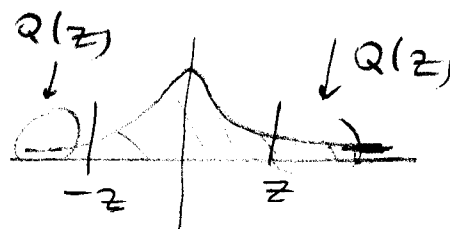
$$= P\{\bar{P}_r(d_{\text{max}})(\text{dB}) + X_\sigma > \gamma\}$$

$$= P\{X_\sigma > \gamma - \bar{P}_r(d_{\text{max}})(\text{dB})\}$$

$$= P\{\sigma Z > \gamma(\text{dB}) - \bar{P}_r(d_{\text{max}})(\text{dB})\}$$

$$= P\{Z > \frac{\gamma(\text{dB}) - \bar{P}_r(d_{\text{max}})(\text{dB})}{\sigma}\}$$

$$= Q\left(\frac{\gamma(\text{dB}) - \bar{P}_r(d_{\text{max}})(\text{dB})}{\sigma}\right)$$



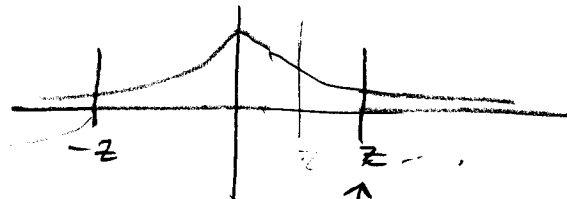
[useful relationships: $Q(-z) = 1 - Q(z)$

(due to symmetry of Gsn. distribution)

$$Q(0) = 1/2 .]$$

- See Table F.1, on p. 647, Rappaport.

- Note that only $Q(z) \leq 1/2$ values are given.
(i.e. for $z \geq 0$.)



$$Q\left(\frac{\gamma(\text{dB}) - \bar{P}_r(d_{\max})(\text{dB})}{\sigma}\right) = 1 - Q\left(\frac{\bar{P}_r(d_{\max})(\text{dB}) - \gamma(\text{dB})}{\sigma}\right)$$

0.01
↓
requires $z \approx \underline{\underline{2.3}}$
from Table F.1

$$\therefore \bar{P}_r(d_{\max})(\text{dB}) - \gamma(\text{dB}) = 2.3\sigma$$

$$\gamma(\text{dB}) = \bar{P}_r(d_{\max})(\text{dB}) - 2.3\sigma$$

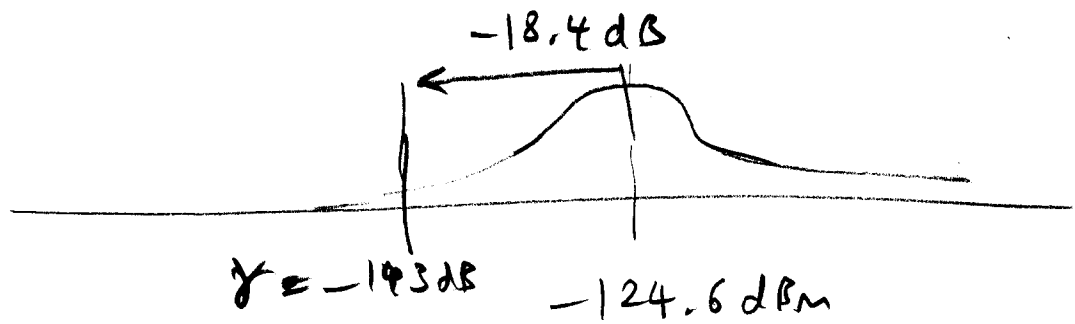
$$= 10 \log_{10}(0.35 \times 10^{-12} \text{ mW}) - 2.3 \cdot (8)$$

$$= -124.6 \text{ dBm} - 18.4 \text{ dB}$$

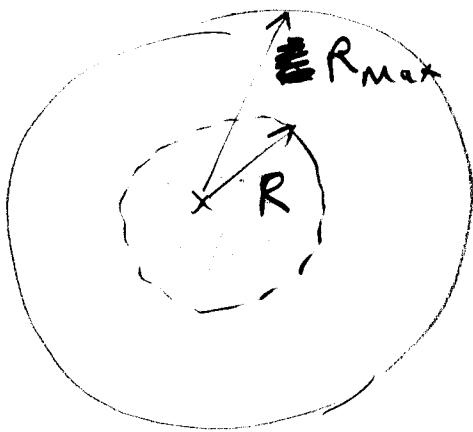
$$= \boxed{-143 \text{ dBm}}$$

↓
the ^{maximum} threshold required (for $P_{\text{out}} = 0.01$)

i.e.,



- The analysis above was done for the worst-case user. This may be too conservative. A more realistic analysis involves the determination of the coverage area based on $V(\gamma) \equiv$ the percentage of useful service area, ($\equiv$ the percentage of the area with a received signal $\geq \gamma$.)



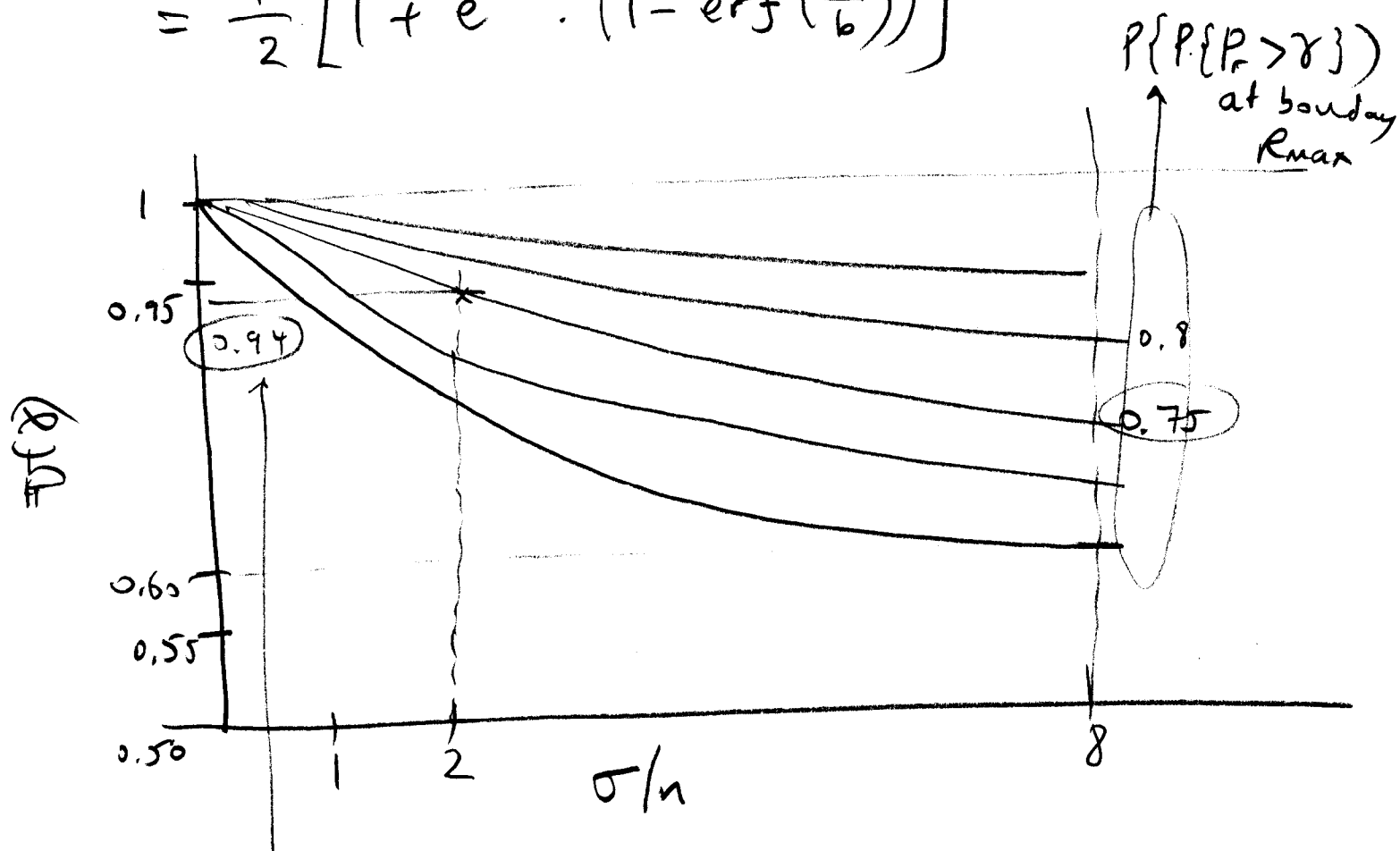
$$\begin{aligned}
 V(\gamma) &\equiv P\{\text{the received signal is } \geq \gamma\} \\
 &\equiv \int_0^{R_{\max}} P\{\text{received signal } \geq \gamma \mid \text{mobile user is at a distance } r \text{ from B.S.}\} \cdot f_R(r) dr \\
 &\quad \uparrow \\
 &\quad \text{condition on } r
 \end{aligned}$$

$$= \frac{1}{\pi R_{\max}^2} \int_0^{2\pi} \int_0^{R_{\max}} \mathbb{I}\{P_r(r) > \gamma\} r dr d\theta$$

$$P_r(r) > \gamma = Q\left(\frac{\gamma - \bar{P}_r(r)}{\sigma}\right)$$

see Rappaport

$$= \frac{1}{2} \left[1 + e^{1/b^2} \cdot (1 - \operatorname{erf}(\frac{1}{b})) \right]$$



e.g. For $\sigma = 8 \text{ dB}$
 $n = 4$.

$\sigma/n = 2$: 0.75 = boundary $\Rightarrow$ 94% area coverage

Part II: TRAFFIC and QUEUING

- Independent of multiplexing technique, traffic/queuing theory characterizes the efficiency of channel usage.

Assumptions: Assume that:

- ① ¹ The number of users is very large,
² The rate of service requests from each user is small.

Then, the service requests ("arrivals") can be considered to be random, and independent, and approximated by a Poisson process.

- ② The holding times (i.e. how long you are on the phone) is also assumed to be random & independent.
- ③ Assume that the ~~sta~~ traffic process is stationary in time, & assume that it is ergodic.

- Erlang: (a dimensionless unit); Measure of traffic ^(rate) ~~rate~~ carried.

Ex A radio channel that is occupied for 30 minutes during 1 hour carries 0.5 Erlangs of traffic.

(- used mainly for voice ~~data~~ traffic.)

- Trunked radio system:

≡ Each user is allocated a channel on a per-call basis.

(Upon termination, the channel is returned to the pool of available channels.)

- Grade of service: ≡ Measure of the ability of a user to access a trunked system during the busiest hour.

- $\updownarrow$ inversely proportional to

the probability that a call is blocked, (i.e., a call is ~~not~~ placed, but due to 0 # of channels being available, user cannot ~~send~~ make the call.)

- Traffic intensity : (offered by a user)

$$= \text{call request rate} \times \text{Holding time}$$

$$\underbrace{A_u}_{\substack{\text{Erlangs} \\ \uparrow \\ \text{dimensionless}}} = \underbrace{\mu}_{\substack{\uparrow \\ \text{average \# of call requests per second or minute}}} \cdot \underbrace{H}_{\substack{\downarrow \\ \text{(seconds or minutes)}}$$

- System with U users, the total traffic intensity

$$A = U A_u = U \cdot \mu \cdot H$$

(assuming same for each user)

- Traffic intensity per channel:

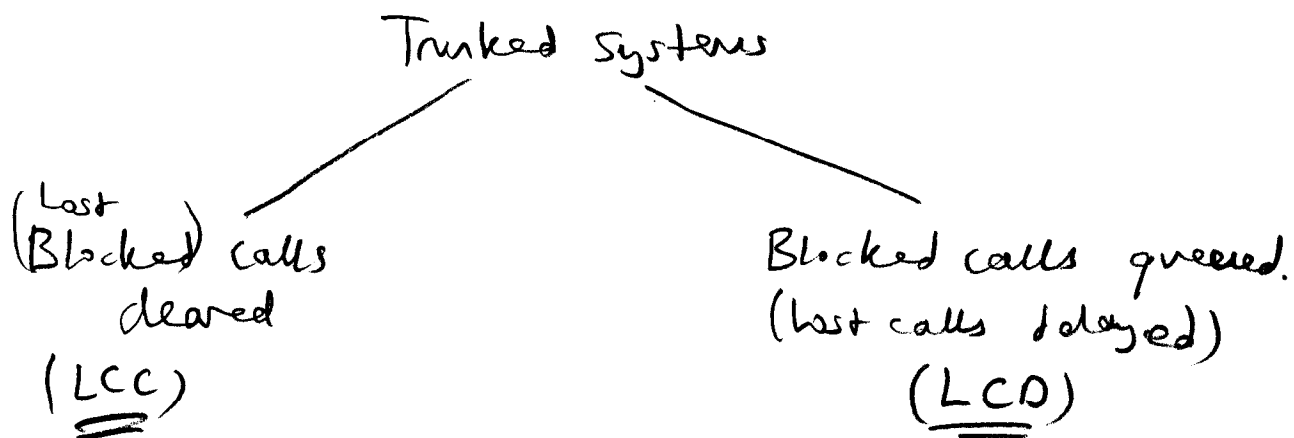
$$A_c = \frac{U A_u}{C}$$

$\uparrow$
of channels

- Maximum ~~the~~ amount of traffic carried by -

System is C Erlangs, (C is the # of channels)

Ex] AMPS (analog, 1G): ^{grado of service} GOS of 2% blocking.
 ↓
 (during busiest hour)



• user may retry the call.

GOS: blocking probability.

GOS: $P\{\text{calls will have delays} > t \text{ seconds}\}$

① LCC systems: (Blocked calls cleared).

- Assume that Service times (holding times) are exponentially distributed.
- Finite # of channels. (C)
- A: total load offered.
- calls arrive $\sim$ Poisson process

C channels

U users

over all users

λ : ^{total mean} call arrival rate (per unit time) [arrival rate]

H : average holdup time.

A : ^{total} offered load.

A_u : ~~avg~~ average load (per user)

$$(A_u = \lambda_1 \cdot H)$$

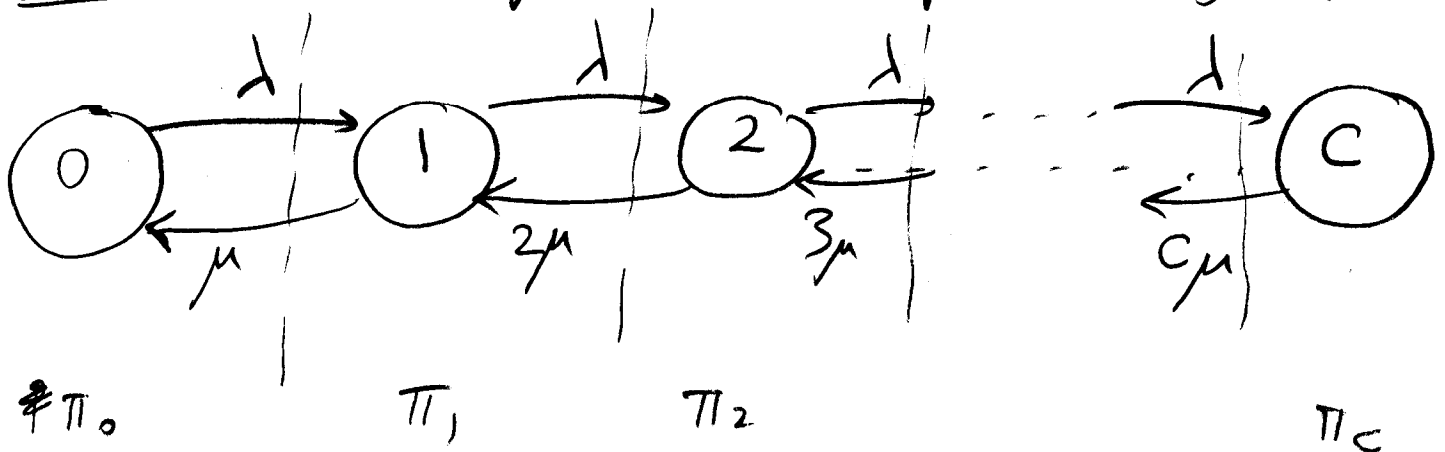
↑ arrival rate of user 1.

$$[A = U A_u = \lambda H]$$

$P\{\text{blocking}\} = P\{\text{none of the channels free for access}\}$

- Markov process:

each state $\longleftrightarrow$ # of channels occupied in system:



Then: $P\{\text{blocking}\} = P\{\text{state } c\}$.

1. Solve $\pi P = \pi$ (steady state probabilities) and look at π_c .

2. or write the balance equations for each node in the flow graph.

$$\lambda \pi_0 = \mu \cdot \pi_1$$

$$\lambda \pi_1 = 2\mu \cdot \pi_2$$

$$\lambda \pi_2 = 3\mu \pi_3$$

⋮

$$\lambda \pi_{i-1} = i\mu \cdot \pi_i$$

$$1 \leq i \leq c$$

$$\rightarrow \pi_1 = \left(\frac{\lambda}{\mu}\right) \pi_0$$

$$\rightarrow \pi_2 = \left(\frac{\lambda}{2\mu}\right) \pi_1 = \frac{\lambda^2}{2\mu^2} \pi_0$$

$$\rightarrow \pi_3 = \left(\frac{\lambda}{3\mu}\right) \pi_2 = \frac{\lambda^3}{3 \cdot 2 \cdot 1 \cdot \mu^3} \pi_0$$

$$\rightarrow \pi_i = \frac{\lambda^i}{i! \mu^i} \pi_0$$

and $\sum_{i=0}^c \pi_i = 1$

Now, $\sum_{i=0}^c \left[\frac{\lambda^i}{i! \mu^i} \right] \pi_0 = 1$.

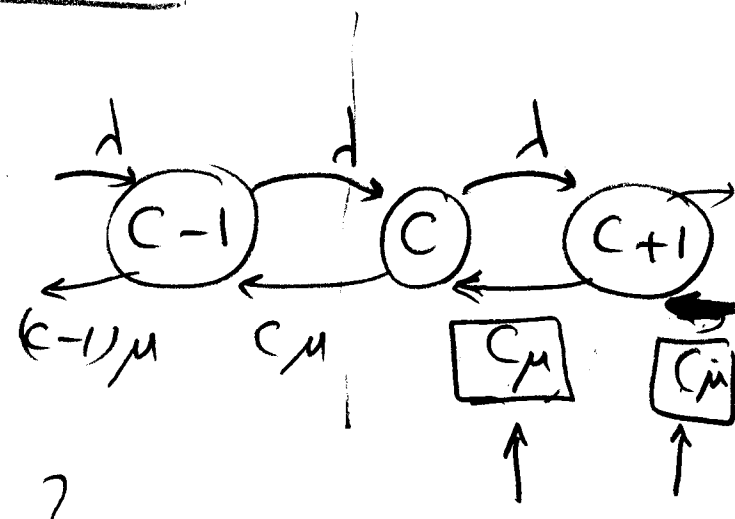
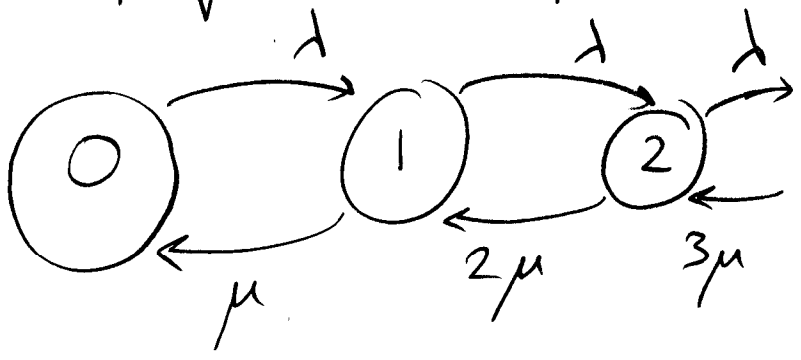
$$\pi_0 = \frac{1}{\sum_{i=0}^c \frac{\lambda^i}{i! \mu^i}}$$

↑ Erlang's formula

$$\boxed{\pi_c} = \frac{\frac{\lambda^c}{c! \mu^c}}{\sum_{i=0}^c \frac{\lambda^i}{i! \mu^i}} = \frac{\left(\frac{\lambda}{\mu}\right)^c \cdot \frac{1}{c!}}{\sum_{i=0}^c \left(\frac{\lambda}{\mu}\right)^i \cdot \frac{1}{i!}}$$

(B) Lost call delayed (LCD discipline)

note: # of channels occupied.



$$\lambda \pi_0 = \mu \pi_1$$

$$\lambda \pi_1 = 2\mu \pi_2$$

$$\lambda \pi_2 = 3\mu \pi_3$$

⋮

$$\lambda \pi_{C-2} = (C-1)\mu \cdot \pi_{C-1}$$

$$\lambda \pi_{C-1} = C\mu \cdot \pi_C$$

$$\left. \begin{array}{l} \lambda \pi_{i-1} = i\mu \cdot \pi_i \\ 1 \leq i \leq C. \end{array} \right\}$$

$$\lambda \pi_C = \mu \cdot C \cdot \pi_{C+1}$$

$$\left. \begin{array}{l} \lambda \pi_i = \mu \cdot C \cdot \pi_{i+1} \\ C \leq i < \infty. \end{array} \right\}$$

$$\sum_{i=0}^{\infty} \pi_i = 1$$

Solving for π_0 : $\pi_0 = \frac{1}{\sum_{i=0}^{C-1} \left(\frac{\lambda}{\mu}\right)^i \frac{1}{i!} + \frac{1}{C!} \left(\frac{\lambda}{\mu}\right)^C \left(\frac{1}{1 - \frac{\lambda}{\mu C}}\right)}$

$$\pi_i = \begin{cases} \left(\frac{\lambda}{\mu}\right)^i \frac{1}{i!} \pi_0 & 0 \leq i \leq c \\ \frac{1}{c!} \left(\frac{\lambda}{\mu}\right)^i \frac{1}{c^{i-c}} \pi_0 & i \geq c \end{cases}$$